

The Future of Real-time Big Data:

# In-Memory Data Grids

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#### **Presenter**





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## "Intellectuals solve problems. Geniuses prevent them." -- Albert Einstein

- 20+ years of shipping scalable, distributed enterprise software
- Vendor neutral. Merit-based. Biased towards simplicity/elegance.
- Author, mentor & conference speaker
- Founder and past persident of New England WebLogic User Group
- M.S. in Computer Science, Adjunct Professor at Northeastern University



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## Agenda



**Definitions:** 

What is Big Data? What is Real-time Big Data Analytics? What is High-velocity data stream computing

**Business Drivers** 

Big Data Universe: Hadoop, NoSQL and Distributed Data Grids

**Distributed Data Grids:** 

What are Distributed Data Grids Charateristics of a Distributed Data Grid

Major players, and the winner is....

From Technlogies to Platforms

The Ultimate: Unified Big Data processing



- Not an IMDG tutorial
- Not In-depth comparison of IMDGs (things change too fast)
- Data modeling, partitioning strategy (hub/spoke, data affinity, etc) (separate presentation)
- Analysis of competing frameworks suitable for Real-Time B.D.A. (separate presentation)
- CAP Theorem discussion ③



Data sets too large and complex to process using standard DB tools and data processing applications



#### **Real-time** Big Data Analytics





### Focus on Velocity + Variety

#### Rapidly changing data (at rest or in-motion):

**Business activities** Examples: trades, orders

**Examples:** market data, current customer location

#### Operational monitoring

Example: application log entries



#### **Operational Intelligence:**

- Analyze stream of business activities and external stimuli on-the fly
- React to them (preferably) instantaneously
- First-mover competitive advantage (new market conditions, customer wishes, etc)
- Real time data stream processing is critical, otherwise business value is lost.

#### Real-time Big Data Analytics examples:

- Dynamic pricing (e-commerce)
- High-frequency trading
- Network security threats
- Credit card fraud prevention
- Factory floor data collection, RFID
- Mobile infrastructure, machine to machine (M2M) applications
- Prescriptive or Location-based applications
- Real-time dashboards, alerts, and reports



- In-Memory Computing Racing Towards Mainstream Adoption
- Market to Reach \$1 Billion by 2016
- Issues so far:
  - Lack of standards,
  - Scarcity of skills,
  - Relative architectural complexity,
  - Monitoring and management challenges



- Advanced queries or complex transactions on huge datasets
- Get results in (near) real-time

## In-motion:

- Rapidly changing, massive stream(s) of data (events)
- Multiple sources, formats
- Needs to be processed in-flight
- Events persisted or not, depending on value

## At rest:

- Data already persisted, change notification
- Re-ingest, re-process (deltas only?)

Typically, a combination of both



### Hadoop

Spark Shark (Spark on Hive) 40x

Still needs Hadoop: HDFS, YARN, Pig, Hive, HBase, JobTracker. TaskTracker,...

Open-source, batch-oriented

### MPP DBs

- Teradata
- Vertica
- Greenplum

Proprietary, complex, expensive

## **NoSQL DBs**

Column: Cassandra, Hbase, Redshift Document/K-V: MongoDB, CouchDB, Riak Graph: Neo4j, etc.

No tx, no referential integrity, triggers, foreign keys

## **In-Memory Data Grids**

- Coherence (commercial
- Hazelcast (OSS)
- Terracotta (commercial)
- Gemfire (commercial)
- GridGain (commercial)
- Gigaspaces XAP (OSS)

## The Origins of In-Memory Data Grids





- RAM is the new disk
- DISK is the new tape

SQL (not Turing complete) MapReduce (low granularity, invasive), often not enough

- From: Disk IO optimization
- To: Data management over network (data in a grid)

## **The Concept**





#### Data partitioning:

- Logical (on key) % LOGICAL\_PARTITION\_COUNT (Hazelcast = 271)
- Physical (one or more logical)
- Partition ownership, replicas reassigned on membership change
- Can determine physical location by key, to dispatch logic



## Hadoop family

Spark/Shark

### **Apache Storm**



## Traits of a Real-time Analytics platform



#### Low Transactional Latency

In-memory transaction speeds, minimize network trips, response in ms

#### Low Data Latency Real-time data, rather than stale data weeks or months old

- High Throughput, High Scalability Horizontal and vertical
- Stream processing/messaging/CEP "Continuous Query"

#### Resilience

Stateless, decentralized, peer-peer - no single point of failure, Partitioned with backup replicas. nodes can fail without data loss and join without any data corruption

- Elastic, simple cluster management No"Mesos on top of ZooKeeper, etc."
- Programmatic, performant SQL-like querying (data affinity)
- Txs , including XA (when needed)

## Major players - comparison



Feature	Hazelcast 3.2	GemFire 7	Oracle Coherence 12.1.3
Distributed data/collections	DistributedMap MultiMap Dist. Queue Dist. Events	No Multimap	No Multimap
Distributed Concurrency	Dist. Lock Dist. Atomic Ref Dist. Atomic Dist Semaphore	No Dist. Atomic Ref. No Dist. Semaphore	No Dist. Atomic Ref. No Dist. Semaphore
Distributed Computing	Executor Service, Key Affinity, "Map/Reduce"	Convoluted, Invasive	Convoluted, Invasive
Resilience	In-memory replicas Recovery WAN replication	Yes	Yes
Elasticity/Cluster Management	Yes	Yes	Yes
Transactions	Local + XA	Local + XA	Local + XA
Distributed Querying	SQL (subset) Predicate Continuous Query	Yes	Yes
License	Apache 2.0 + Commercial	Commercial	Commercial



## hazelcast

- The most intellectually elegant distributed in-memory data/compute grid.
- Minimalism as design aesthetic: Non-intrusive, No dependencies 2.6MB single jar library.
- Implements Java APIs (Map, List, Set, Queue, Lock) in a distributed manner
- Distributed Execution Framework (extension of Java's Executor Service) Distributed Queries (SQL/Predicate), Data affinity (execution on specific node execution)
- Cluster management Java API included (dead simple) Auto-discovery of nodes and intelligent synchronization
- Apache License 2, commercial extensions + support



**import** java.util.concurrent.Callable;

```
public class Echo implements Callable < String >, Serializable {
```

```
String input = null;
```

```
public Echo() { }
public Echo(String input) {
    this.input = input;
}
public String call() {
    return Hazelcast.getCluster().getLocalMember().toString() + input;
}
```

```
ExecutorService executorService = Executors.newSingleThreadExecutor();
Future<String> future = executorService.submit (new Echo("myinput"));
```

```
String result = future.get();
```

...



import com.hazelcast.core.lExecutorService; import java.util.concurrent.Callable; import java.util.concurrent.Future; import java.util.Set;

public void echoSomewhere(String input, Member member) {
 Callable<String> task = new Echo(input);
 HazelcastInstance hz = Hazelcast.newHazelcastInstance();
 IExecutorService executorService = hz.getExecutorService("default");

Future<String> future = executorService.submitToMember(task, member); Future<String> future = executorService.submitToKeyOwner(task, key); Map<Member, Future<String>> futures = executorService.submitToMembers(**new** Echo(input), members);

String echoResult = future.get();

Execution callbacks, cancellation APIs provided



#### Code example: Data Affinity(1)

```
HazelcastInstance instance = Hazelcast.newHazelcastInstance(cfg);
Map mapa = instance.getMap("mapa");
Map mapb = instance.getMap("mapb");
Map mapc = instance.getMap("mapc");
```

```
// different entries, but the operation will take place on the same member
mapa.put("key1", value);
mapb.get("key1");
mapc.remove("key1");
```

// lock operation will still execute on the same member of the cluster
instance.getLock ("key1").lock();

// consistent with distributed execution
instance.getExecutorService().executeOnKeyOwner(runnable, "key1");





#### public class OrderKey implements Serializable, PartitionAware {

```
int customerId;
int orderId;
public OrderKey(int orderId, int customerId) {...}
public Object getPartitionKey() { return customerId; }
```

}

```
Map mapCustomers = instance.getMap("customers");
Map mapOrders = instance.getMap("orders");
mapCustomers.put(1, customer);
mapOrders.put(new OrderKey(21, 1), order);
```

## public static class OrderDeletionTask implements Callable<Integer>, PartitionAware, Serializable { .... }

ExecutorService es = instance.getExecutorService(); OrderDeletionTask task = **new** OrderDeletionTask(customerId, orderId); Future future = es.submit(task); **int** remainingOrders = future.get();



```
public class Employee implements Serializable {
    private String name;
    private int age;
    private boolean active;
    private double salary;
}
```

HazelcastInstance hz = Hazelcast.newHazelcastInstance(cfg); IMap map = hz.getMap("employeeMap");

```
Set<Employee> employees = (Set<Employee>) map.values(
new SqlPredicate("active AND age < 30"));
```

#### // or JPA-like criteria

Predicate predicate = e.is("active").and(e.get("age").lessThan(30)); Set<Employee> employees = (Set<Employee>) map.values(predicate);



```
HazelcastInstance hz = Hazelcast.newHazelcastInstance(cfg);
Cluster cluster = hz.getCluster();
cluster.addMembershipListener(new MembershipListener(){
public void memberAdded(MembershipEvent membersipEvent) {
System.out.println("MemberAdded " + membersipEvent);
}
```

Member localMember = cluster.getLocalMember(); System.out.println ("my inetAddress= " + localMember.getInetAddress()); Set setMembers = cluster.getMembers(); for (Member member : setMembers) { System.out.println ("isLocalMember " + member.localMember()); System.out.println ("member.inetaddress " + member.getInetAddress()); System.out.println ("member.port " + member.getPort());

#### Example: Admin console

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Scalable Computing:

- DB Caching
- Web/Session Clustering
- High-speed K-V Data Store
- Distributed queues, topics
- (Fast) SOA Services



- RAM currently maxes out at ~ 640GB/server
- Still more expensive than SSDs and HDDs
- Garbage collection

- Cost and capacity limitations will disappear over time
- Off-heap memory and specialized JVMs (Azul, etc.)



In-Memory Data/Compute Grids not enough:

- Backing DB(s)
- Enterprise Message Store(s)

Example: GemFire GreenPlum DB



- Multi-Source Data Harvesting, Ingestion, Transformation
- Data-type agnostic:

Highly Structured (RDBMS data) Semi-structured (documents, key-value tuples) In-Motion and At Rest

- Distributed Computing, not just Data Tx and non-Tx Batch and/or Real-time Processing
- The fewer moving parts the better

Data Warehousing -> Hadoop/MapReduce -> Unified Big Data ProcessingRelational + BatchBatchAll of the above









- One order of magnitude faster
- Data And Logic colocated, affinity-able (with careful design)
- Flexibility:
  - Turing complete
  - SQL/Predicate queries
  - MapReduce framework on top
- Scalable and Elastic, simply (automatic repartitioning)
- Current limitations (cost, GC) fading fast





## http://www.numatica.com

# Appendix

## Berkeley Big-data Analytics Stack (BDAS)



Applications built over BDAS